

# **Family Example: Some Facts**

- **respects(barb,dan).**
- **respects(barb,katie).**
- **respects(dan,brian).**
- **respects(dan,barbara).**

# Some Queries

**?- respects(barb,katie).**

**yes**

**?- respects(barb,X).**

**X=dan; % you type ; and RETURN**

**X=katie; % you type ; and RETURN**

**no**

**% comment to end of line**

# Some More Queries

**?- respects(X,Y).**

**X=barb Y=dan**

**?- respects(barb,katie),respects(katie,barb).**

**no            % conjunctive query: , = AND**

**?- respects(barb,X),respects(X,barb).**

**X=dan**

# Some Rules

- **mutualRespect(X,Y):- respect(X,Y),respect(Y,X).**
- **selfRespect(X,X):- respect(X,X).**
  - **% selfRespect(X,Y):- respect(X,Y), X=Y.**
- **sister(X,Y):- female(X),parents(X,Ma,Pa),  
parents(Y,Ma,Pa),X\==Y.**

# Some More Rules

- **ancestor(Person,Anc) :- parent(Person,Anc).**
- **ancestor(Person,Anc):- parent(Person,X),ancestor(X,Anc).**
  
- **parent(a,b).**
- **parent(b,c).**
- **parent(a,e).**
- **parent(b,f).**

**?- ancestor(a,X).**

**% returns in order b, e, c, f**

# Script: Running Prolog

CREATE FILE NAMED ancestor WITH RULES AND FACTS

> script scriptFileName (CCC Alpha UNIX: garden,wpi,ccc,cpu,...)

% pl TO ENTER PROLOG

Welcome to SWI-Prolog (Version 2.9.9)

? - **consult('ancestor').** READS IN FILE NAME ancestor

? - ancestor(a,X).

X = b ;

X = e ;

X = c ;

X = f ;

No

? - **halt.** TO EXIT PROLOG

% **exit** TO EXIT SCRIPT

# Unification:

## Pattern Matching with Variable Binding

- same predicate
- same number of terms
- for each term:
  - same constant
  - variable and constant
    - variable bound to constant
  - variable and variable
    - if both bound, must be to same
    - if not, bound to each other

# The Prolog Inference Engine

- **schoolmates(X,Y) :-**  
    **student(X, EnterX, ExitX),**  
    **student(Y, EnterY, ExitY),**  
    **X\==Y,**  
    **overlap(EnterX, ExitX, EnterY, ExitY).**
- **overlap(EnterX,ExitX,EnterY,ExitY) :-**  
    **EnterY >= EnterX,**  
    **EnterY < ExitX.**
- **overlap(EnterX, ExitX, EnterY, ExitY) :-**  
    **EnterX >= EnterY,**  
    **EnterX < ExitY.**
- **student(bill, 1992,1996).**
- **student(sue,1986,1990).**
- **student(stu, 1994, 1998).**



- ?- schoolmates(bill,X).
- **UNIFICATION** with schoolmates rule head
- **Goals to Prove:**
  - student(X, EnterX, ExitX),
  - student(Y, EnterY, ExitY),
  - $X \text{ \textbackslash } == Y$ ,
  - overlap(EnterX, ExitX, EnterY, ExitY).
- **UNIFICATION** of first goal with student(bill, 1992, 1996).
- **Goals to Prove:**
  - student(bill, 1992, 1996),
  - student(Y, EnterY, ExitY),
  - $\text{bill} \text{ \textbackslash } == Y$ ,
  - overlap(1992, 1996, EnterY, ExitY).

# Execution Trace.

?- trace.

?- schoolmates(X,Y).

**Call: ( 7) schoolmates(\_G165, \_G166)**

**Call: ( 8) student(\_G165, \_L131, \_L132)**

**Exit: ( 8) student(bill, 1992, 1996)**

**Call: ( 8) student(\_G166, \_L133, \_L134)**

**Exit: ( 8) student(bill, 1992, 1996)**

**Call: ( 8) bill\==bill Fail: ( 8) bill\==bill**

# Trace Continuation 1

**Redo: ( 8) student(\_G166, \_L133, \_L134)**

**Exit: ( 8) student(sue, 1986, 1990)**

**Call: ( 8) bill\==sue Exit: ( 8) bill\==sue**

**Call: ( 8) overlap(1992, 1996, 1986, 1990)**

**Call: ( 9) 1986>=1992**

**Fail: ( 9) 1986>=1992**

# Trace Continuation 2

- **Redo: ( 8) overlap(1992, 1996, 1986, 1990)**
- **Call: ( 9) 1992>=1986**
- **Exit: ( 9) 1992>=1986**
- **Call: ( 9) 1992?**
- **Fail: ( 9) 1992?**
- **Fail: ( 8) overlap(1992, 1996, 1986, 1990)**

# Trace Continuation 3

**Redo: ( 8) student(\_G166, \_L133, \_L134)**

**Exit: ( 8) student(stu, 1994, 1998)**

**Call: ( 8) bill\==stu Exit: ( 8) bill\==stu**

**Call: ( 8) overlap(1992, 1996, 1994, 1998)**

**Call: ( 9) 1994>=1992**

**Exit: ( 9) 1994>=1992**

**Call: ( 9) 1994? ?**

**Exit: ( 9) 1994? ?**

**Exit: ( 8) overlap(1992, 1996, 1994, 1998)**

**Exit: ( 7) schoolmates(bill, stu)**

# Results

?- schoolmates(X,Y).

**X = bill            Y = stu;**

**X = stu            Y = bill;**

**No**

**?- listing.**

**student(bill, 1992, 1996).**

**student(sue, 1986, 1990).**

**student(stu, 1994, 1998).**

**schoolmates(A, B) :-**

**student(A, C, D),**

**student(B, E, F),**

**A\==B,**

**overlap(C, D, E, F).**

**overlap(A, B, C, D) :- C>=A, C<B.**

**overlap(A, B, C, D) :- A>=C, A<D.**

# Prolog: Built-in Types

- **atoms**
- **numbers**
- **record-like structure: functor(ListOfComponents)**
- **lists**



# Lists

**delimiters: [ ]**

**separator: ,**

**[good, bad, ugly]**

**empty list: [ ]**

**head/car, tail/cdr: [Head|Tail]**

**member( X, [ X | \_ ] ) .**

**member( X, [ \_ | Rest ] ) :- member( X, Rest ) .**

# Example: APPEND

## Invertibility

### No designated input arguments and return value

**append( [], Whole ,Whole ).**

**append( [HStart|TStart], End, [HStart|TWhole] ) :-  
append( TStart, End, TWhole ).**

**?-append([good],[bad,ugly], Movie).**

**?-append([good,bad],What,[good,bad,ugly]).**

**?-append(What,[ugly],[good,bad,ugly]).**

**?-append(What,WhatElse,[good,bad,ugly]).**

# Example: MYSORT

- **mysort(Xs,Ys) :- permutation(Xs, Ys), ordered(Ys).**
- **ordered([X]).**
- **ordered([X, Y|Ys]) :- X =< Y, ordered([Y|Ys]).**
- **permutation([], []).**
- **permutation(Xs, [Z|Zs]) :- remove(Z, Xs, Ys),  
permutation(Ys, Zs) .**
- **remove(X, [X|Xs], Xs).**
- **remove(X, [Y|Ys], [Y|Zs]) :- remove(X, Ys, Zs).**

# More Example Rules

- **count\_up([], 0).**
- **count\_up([X|R], Count):-  
    count\_up(R, Subcount),  
    Count is 1 + Subcount.**
- **min([H|T],Z):- minsofar(T, H, Z).**
- **minsofar([], X, X).**
- **minsofar([H|T], X, Z) :- X =< H, minsofar(T, X, Z).**
- **minsofar([H|T], X, Z) :- H < X, minsofar(T,H, Z).**

# bagof

**commitment(1,1,1,97,9,1).**

**commitment(2,1,1,98,9,1).**

**commitment(3,1,2,97,9,1).**

**commitment(4,2,1,97,9,1).**

**commitment(5,1,1,97,10,1).**

**commitment(6,1,1,97,8,2).**

**commitment(7,1,2,97,7,4).**

**in97(Id) :- commitment(Id,\_,\_,97,\_,\_).**

**count97commits(C) :-**

**bagof(X, in97(X), L),**

**count\_up(L, C).**

```
| ?- ['bagofadvice'].    %OR consult('bagofadvice').  
| ?- in97(Id).  
Id = 1 ;  
Id = 3 ;  
Id = 4 ;  
Id = 5 ;  
Id = 6 ;  
Id = 7 ;  
no  
| ?- bagof(X,in97(X),L).  
X = _0  
L = [1,3,4,5,6,7] ;  
no | ?- count97commits(C).  
C = 6
```

# Resolution Order

- **tries to unify** in fixed order
  - top-down** in list of facts and rules
- **tries to satisfy** (sub)goals in fixed order
  - left to right** in RHS of rule
- **depth-first** search
  - new goals put a front of list of goals to solve**

# Order of Rules Matters Order of Subgoals Matters

**ancestor(Person, Anc) :- parent(Person, Anc).**

**ancestor(Person, Anc):- parent(Person, X), ancestor(X, Anc).**

**parent(a, b).**

**parent(b, c).**

**parent(a, e).**

**parent(b, f).**

**?- ancestor(a, X).**

**% returns in order b, e, c, f**



# Order of Rules Matters Order of Subgoals Matters

**ancestor(Person,Anc) :- parent(Person, X), ancestor(X, Anc).**

**ancestor(Person,Anc) :- parent(Person, Anc).**

**parent(a,b).**

**parent(b,c).**

**parent(a,e).**

**parent(b,f).**

**?- ancestor(a,X).**

**% returns in order c, f, b, e**

# Order of Rules Matters Order of Subgoals Matters

**ancestor(Person, Anc):- ancestor(X,Anc), parent(Person,X).**

**ancestor(Person, Anc) :- parent(Person, Anc).**

**parent(a,b).**

**parent(b,c).**

**parent(a,e).**

**parent(b,f).**

**?- ancestor(a,X).**

**% returns no answer, runs out of heap**

# SOUND, but NOT COMPLETE

- A set of inference rules or an inference procedure is SOUND if everything(theorem) that is derived using those rules or that procedure logically follows from the facts (and rules).
- A set of inference rules or an inference algorithm is COMPLETE if everything(theorem) that logically follows from the facts (and rules) can be derived using those rules or that procedure.
- Prolog is SOUND, but NOT COMPLETE

# (Extralogical) Cut !

- **! can appear in RHS of rule (clause body)**

`a :- b, c ,!, d, e.`

- **! is a goal which always succeeds, exactly once**
- **no backtracking through !**
- **no possibility to use additional clauses to prove goal**

`insertIfNotThere(X, L, L) :- member(X, L) , ! .`

`insertIfNotThere(X, L, [X|L]).`

- **! can save computation time, BUT at a cost**

- **!** can be used to imitate if-then-else

- **max(X, Y, X) :- X >= Y.**

- **max(X, Y, Y).**

**?-max(5, 4, M).**

**M=5;**

**M=4**

- **max(X, Y, X) :- X >= Y, ! .**

- **max(X, Y, Y).**

**?-max(5, 4, M).**

**M=5;**

**No**

# Cut ! CONTINUED

- **BUT  $\text{max}(X, Y, Y)$ . is not a true fact.**
- **It can only be understood in context of previous statement, i.e. must mentally execute to understand, and thus loses advantage of declarative programming**
- **$\text{max}(X, Y, X) :- X \geq Y$ .**
- **$\text{max}(X, Y, Y) :- X < Y$ .**
  - easier to understand (& parallelize), longer to execute

# Negation in Prolog

- **not in Prolog is not equivalent to logical not**
- **not can only appear in RHS of rule( clause body)**
- **Negation as Failure**
  - `not(X) :- X, !, fail.`
  - `not(_).`
- **Closed World Assumption**

**All relevant knowledge (facts and rules) about domain are included in knowledge base**

- (if so, if you can't prove a predicate **X**, you could conclude **not X**)

# Negation in Prolog

- When you ‘prove’ something, adding additional info should not affect it.
- Using Negation as Failure, this is not the case.

parent(bob, amy). % only thing in Knowledge Base

?- not(mother(bob, amy)).

Yes % since no rule defining predicate mother

female(amy). % add to KB

mother(X,Y) :- parent(X,Y), female(Y). % add to KB

?- not(mother(bob, amy)).

No

- nonmonotonic reasoning



# Example: LAST

- **member(X, [X|\_]).**
- **member(X, [\_|Ys]) :- member(X, Ys).**
- **last([X|[ ]], X).**
- **last([X|Y], Z) :- last(Y, Z).**

# Example: COUNT\_UP

- **count\_up([ ],0).**
- **count\_up([X|R],Count) :-  
    count\_up(R, Subcount),  
    Count is 1 + Subcount.**
- **sum\_up([ ],0).**
- **sum\_up([X|R],Total) :-  
    sum\_up(R, Subtotal),  
    Total is X + Subtotal.**

# **Example: AVERAGE**

**average(L, Average):-**

**sum\_up(L, Total),**

**count\_up(L, Count),**

**Average is Total/Count.**

# **Example: PAYROLL**

**hours(emp1,10).**

**hours(emp2,20).**

**hours(emp3,30).**

**rank(emp1,a).**

**rank(emp2,a).**

**rank(emp3,b).**

**payscale(a,25).**

**payscale(b,50).**

# PAYROLL (CONTINUED)

- **hours(emp4,40).**
- **rank(emp4,c).**
- **payscale(c,100).**
- **pay(X,Y) :-**
  - hours(X, H),**
  - rank(X, R),**
  - payscale(R, S),**
  - Y is H \* S.**
- **payroll(P) :-**
  - bagof(Y, X^(pay(X,Y)), L),**
  - sum\_up(L, P).**

# The Last Example

**cpi(1,100).**

**cpi(2, 20).**

**cpi(3, 300).**

**cpi(4, 40).**

**cpi(5, 50).**

**taxable(1).**

**taxable(4).**

**taxable(5).**

# The Last Example (CONTINUED)

- **itemtotal(I,T) :-**  
    **cpi(I,Q),**  
    **quantity(I,N),**  
    **taxable(I),**  
    **UT is Q \* N,**  
    **T is UT \* 1.05.**
- **itemtotal(I,T) :-**  
    **cpi(I,Q),**  
    **quantity(I,N),**  
    **not(taxable(I)), T is Q \* N.**
- **totalbill(G) :-**  
    **qreadin(0),**  
    **bagof(T,I^(itemtotal(I,T)),L),**  
    **sum\_up(L,G).**

- **totalbill(G) :-**  
    **qreadin(0),**  
    **bagof(T,I^(itemtotal(I,T)),L),**  
    **sum\_up(L,G).**
- **qreadin(Num) :-**  
    **write( 'Next item, please: '),**  
    **read(X),**  
    **processb(X,Num).**
- **processb(stop,Num):-!.**
- **processb(X,Num) :-**  
    **NN is Num + 1,**  
    **assert(quantity(NN,X)),**  
    **qreadin(NN).**

# **The Last Example (CONTINUED)**



# **Prolog Search**

# Implementing Search in Prolog

- How to represent the problem
- Uninformed Search
  - depth first
  - breadth first
  - iterative deepening search
- Informed Search
  - Hill climbing
  - Graph Search
    - which can do depth first, breadth first, best first, Algorithm A, Algorithm A\*, etc.

# Representing the Problem

- Represent the problem space in terms of two predicates:
  - goal/1
  - arc/3
- **goal(S)** is true iff S is a goal state.
- **arc(S1,S2,N)** is true iff there is an operator of cost N that will take us from state S1 to state S2.
- **arc(S1,S2) :- arc(S1,S2,\_)**

# Eight Puzzle Example

- Represent a state as a list of the eight tiles and o for blank.
- E.g., [1,2,3,4,o,5,6,7,8] for

1	2	3
4	o	5
6	7	8

```
goal ([ 1, 2, 3,  
        4, o, 5,  
        6, 7, 8 ]).
```

```
arc ([ o, B, C,  
        D, E, F,  
        G, H, I ],  
     [ B, o, C,  
        D, E, F,  
        G, H, I ]).
```

# Missionaries and Cannibals

```
% Represent a state as
% [ML,CL,MR,CR,B]
start([3,3,0,0,left]).
goal([0,0,3,3,X]).

arc([ML,CL,MR,CR,left],
    [ML2,CL,MR2,CR,right]):-
    % one M & one C row right
    MR2 is MR+1,
    ML2 is ML-1,
    CR2 is CR+1,
    CL2 is CL-1,
    legal(ML2,CL2,MR2,CR2).

arc([ML,CL,MR,CR,left],
    [ML2,CL,MR2,CR,right]):-
    % two Ms row right
    MR2 is MR+2,
    ML2 is ML-2,
    legal(ML2,CL2,MR2,CR2).

Legal(ML,CL,MR,CR) :-
    % is this state a legal
    one?
    ML>0, CL>0, MR>0, CR>0,
    ML>=CL, MR>=CR.
```

# Depth First Search

```
%% this is surely the simplest  
%% possible DFS.
```

```
dfs(S,[S]) :- goal(S).
```

```
dfs(S,[S|Rest]) :-
```

```
    arc(S,S2),
```

```
    dfs(S2,Rest).
```

# Depth First Search which avoids loops

```
% this version of DFS keeps track of the path as  
% it explores, enabling it to avoid loops. It also  
% returns the path from Start to Goal
```

```
:- ensure_loaded(library(lists)).
```

```
dfs(S,Path) :- dfs1(S,[S],Path).
```

```
dfs1(S,Path,ReversePath) :-  
    goal(S),  
    reverse(Path,ReversePath).
```

```
dfs1(S,SoFar,Path) :-  
    arc(S,S2),  
    \+ (member(S2,SoFar)),  
    dfs1(S2,[S2|SoFar],Path).
```

# Breadth First Search

```
:- use_module(library(queues)).
```

```
bfs(S,Path) :-  
    empty_queue(Q1),  
    queue_head([S],Q1,Q2),  
    bfs1(Q2,Path).
```

```
bfs1(Q,[G,S|Tail]) :-  
    queue_head([S|Tail],_,Q),  
    arc(S,G), goal(G).
```

```
bfs1(Q1,Solution) :-  
    queue_head([S|Tail],Q2,Q1),  
    findall([Succ,S|Tail],  
           (arc(S,Succ), \+member(Succ,Tail)),  
           NewPaths),  
    queue_last_list(NewPaths,Q2,Q3),  
    bfs1(Q3,Solution).
```



# Note on Queues

- `:- use_module(library(queues))`
- **`empty_queue(?Q)`**
  - Is true if Queue has no elements.
- **`queue_head(?Head, ?Q1, ?Q2)`**
  - Q1 and Q2 are the same queues except that Q2 has Head inserted in the front. Can be used to insert or delete from the head of a Queue.
- **`queue_last(?Last, ?Q1, ?Q2)`**
  - Q2 is like Q1 but have Last as the last element in the queue. Can be used to insert or delete from the end of a Queue.
- **`list_queue(+List, ?Q)`**
  - Q is the queue representation of the elements in list List.
- Note: Queues are represented as a pair (L,Hole) where list L ends with a variable unified with Hole.

# Iterative Deepening

```
id(S,Path) :-                                % from(-Var,+Val,+Inc)
    from(Limit,1,5),                          % instantiates Var to #s
    id1(S,0,Limit,Path).                     % beginning with Val &
                                              % incrementing by Inc.
```

```
id1(S,Depth,Limit,[S]) :-                   from(X,X,Inc).
    Depth<Limit,                             from(X,N,Inc) :-
    goal(S).                                  N2 is N+Inc,
                                              from(X,N2,Inc).
```

```
id1(S,Depth,Limit,[S|Rest]) :-
    Depth<Limit,
    Depth2 is Depth+1,
    arc(S,S2),
    id1(S2,Depth2,Limit,Rest).

    / ?- from(X,0,5).
    X = 0 ? ;
    X = 5 ? ;
    X = 10 ? ;
    X = 15 ? ;
    X = 20 ? ;
    X = 25 ? ; ...
    yes
    / ?-
```

# Informed Search

- Hill climbing
- General graph search which can be used for
  - depth first search
  - breadth first search
  - best first search
  - Algorithm A
  - Algorithm A\*

# Hill Climbing

```
hc(Path) :- start(S), hc(S,Path).
```

```
hc(S,[S]) :- goal(S), !.
```

```
hc(S,[S|Path]) :-
```

```
    h(S,H),
```

```
    findall(HSS-SS,
```

```
        (arc(S,SS),h(SS,HSS)),
```

```
        L),
```

```
    keysort(L,[BestH-BestSS|_]),
```

```
    H>BestH -> hc(BestSS,Path)
```

```
        ;(debug("Local max:~p~n", [S]), fail).
```

# Graph Search

The graph is represented by a collection of facts of the form:

**node(S,Parent,Arcs,G,H)** where

- **S** is a term representing a state in the graph.
- **Parent** is a term representing *S*'s immediate parent on the best known path from an initial state to *S*.
- **Arcs** is either *nil* (no arcs recorded, i.e. *S* is in the set open) or a list of terms *C-S2* which represents an arc from *S* to *S2* of cost *C*.
- **G** is the cost of the best known path from the state state to *S*.
- **H** is the heuristic estimate of the cost of the best path from *S* to the nearest goal state.

# Graph Search

In order to use `gs`, you must define the following predicates:

- **goal(S)** true if `S` is a term which represents the goal state.
- **arc(S1,S2,C)** true iff there is an arc from state `S1` to `S2` with cost `C`.
- **h(S,H)** is the heuristic function as defined above.
- **f(G,H,F)** `F` is the metric used to select which nodes to expand next. `G` and `H` are as defined above. Default is *"f(G,H,F) :- F is G+H."*
- **start(S)** (optional) `S` is the state to start searching from.

```

gs(Start,Solution) :-
  retractall(node(____)),
  addState(Start,Start,[],[]),
  gSearch(Path),
  reverse(Path,Solution).

gSearch(Solution) :-
  select(State),
  (goal(State)
   => collect_path(State,Solution)
   | (expand(State),gSearch(Solution))).

select(State) :-
  % find open state with minimal F value.
  findall(f+S,
    (node(S,P,nil,GH),f(G,H,F)),OpenList),
  keysort(OpenList,[%State|Fest]).

expand(State) :-
  debug("Expanding state `p.`n",[State]),
  retract(node(State,Parent,nil,GH)),
  findall(ArcObj-Kid,
    (arc(State,Kid,ArcObj),
     add_arc(State,Kid,G,ArcObj)),
    Arcs),
  assert(node(State,Parent,Arcs,GH)).

```

```

add_arc(Parent,Child,ParentG,ArcObj) :-
  %Child is a new state, add to the graph.
  (!node(Child,____)),
  G is ParentG,ArcObj,
  h(Child,H),
  debug("Adding state `p`with parent `p` and
  cost `p.`n",[Parent,Child,G]),
  assert(node(Child,Parent,nil,GH)),!.

```

```

add_arc(Parent,Child,ParentG,ArcObj) :-
  %Child state is already in the graph.
  %update cost if the new path better.
  node(Child,CurrentParent,Arcs,CurrentGH),
  NewG is ParentG,ArcObj,
  CurrentGNewG,!,
  debug("Updating `p`'s cost thru `p` to
  `p.`n",[State,Parent,NewG]),
  retract(node(Child,____)),
  assert(node(Child,Parent,Arcs,NewGH)),
  %better way to get to any grandkids?
  foreach(member(ArcObj-Child,Arcs),
    (NewObjToChild is NewG,ArcObj,
     update(Child,State,NewObjToChild))).

```

```

add_arc(____).

```

```

collect_path(Start,[Start]) :-
  node(Start,Start,Arcs,0,_H).
collect_path(S,[S|Path]) :-
  node(S,Parent,____),
  collect_path(Parent,Path).

```

# Note on Sorting

- `sort(+L1,?L2)`
  - Elements of the list L1 are sorted into the standard order and identical elements are merged, yielding the list L2.  
|?- `sort([f,s,foo(2),3,1],L).`  
L = [1,3,f,s,foo(2)] ?
- `keysort(+L1,?L2)`
  - List L1 must consist of items of the form *Key-Value*. These items are sorted into order w.r.t. Key, yielding the list L2. No merging takes place.  
|?- `keysort([3-bob,9-mary,4-alex,1-sue],L).`  
L = [1-sue,3-bob,4-alex,9-mary] ?
  - Example:  
youngestPerson(P) :-  
    findall(Age-Person,(person(Person),age(Person,Age)),L),  
    keysort(L,[\_-P]).